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In the talk, we consider two recurrent sequences of random vectors $\{\theta_k\}$ and $\{w_k\}$:

$$\begin{aligned}\theta_{k+1} &= \theta_k + \beta_k(b_1 - A_{11}\theta_k - A_{12}w_k + V_{k+1}), \\ w_{k+1} &= w_k + \gamma_k(b_2 - A_{21}\theta_k - A_{22}w_k + W_{k+1}),\end{aligned}$$

where b_i and A_{ij} are some vectors and some matrices resp., V_{k+1}, W_{k+1} are the noise terms, the scalars β_k, γ_k are step sizes. We consider $\beta_k = c_{0,\beta}(k+k_0)^{-b}$ and $\gamma_k = c_{0,\gamma}(k+k_0)^{-a}$ with exponents a and b satisfying $1/2 < a < b < 1$. These type of sequences $\{\theta_k\}$ and $\{w_k\}$ arise in two-timescale stochastic approximation (TTSA) algorithms as stochastic approximations for the unique solutions (θ^*, w^*) for the system of linear equations

$$A_{11}\theta + A_{12}w = b_1, \quad A_{21}\theta + A_{22}w = b_2.$$

Theorem. *Under the martingale noise setting, let some assumptions are met. Then for all large enough n , one has*

$$\text{dist}_K(\beta_n^{-1/2}\tilde{\theta}_{n+1}, \mathcal{N}(0, \Sigma_\infty^{\text{last}})) \lesssim_{\log n} n^{-1/4},$$

where dist_K is the Kolmogorov distance with respect to all convex sets, $\tilde{\theta}_{n+1}$ is the last iterate.

Similar results are obtained for the sequence $\{w_k\}$, for Polyak-Ruppert averaged TTSA and for Markov noise, although with worse order of approximation error.